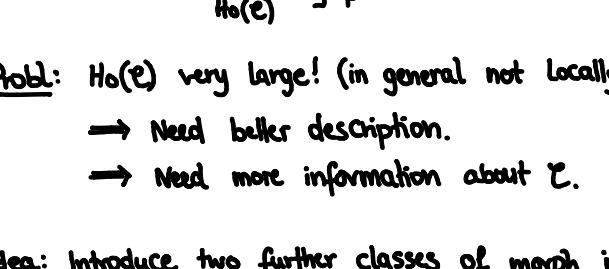


INTRODUCTION TO MODEL CATEGORIES

(main reference: Hovey "Model categories", § 1)

1 Motivation and first definitions

"Homotopy theory" = Category with a notion of "weak equivalence", which is weaker than isomorphism.



Ex: • Topological spaces Top.

• Simplicial sets sSet.

• Positive chain cpx Ch^+ .

• Small categories Cat.

Goal: Understand "objects up to equivalence"

→ Make weak eq. iso by adding formal inverses.
(= model of equivalences)

Constr: (homotopy category) Let $\mathcal{W} \subseteq \text{Mor}(\mathcal{C})$ a class of morphisms.

Define the homotopy cat. $\text{Ho}(\mathcal{C}) = \mathcal{C}[\mathcal{W}^{-1}]$:

• objects: objects of \mathcal{C} .

• morph: $f: X \rightarrow Y$ in \mathcal{C} where \sim weak equivalence

The canonical $\eta: \mathcal{C} \rightarrow \text{Ho}(\mathcal{C})$ satisfies:

• For $w \in \mathcal{W}$, $\eta(w)$ is iso.

• For $F: \mathcal{C} \rightarrow \mathcal{D}$ with Fw iso for all $w \in \mathcal{W}$,

$\eta F: \text{Ho}(\mathcal{C}) \rightarrow \text{Ho}(\mathcal{D})$

and $\text{Ho}(\mathcal{C})$ very large! (in general not locally small)

→ Need better description.

→ Need more information about \mathcal{C} .

Idea: Introduce two further classes of morph in \mathcal{C} :

$\text{Cof} =$ cofibrations ("nice injections")

$\text{Fib} =$ fibrations ("nice surjections")

which behave nicely wrt. equivalences. These will give rise to classes of "weak objects".

• cofibrant objects ($X \rightarrow X$ cofib), η initial obj.

• fibrant objects ($X \rightarrow *$ fib), η terminal obj.

and it will be easier to understand them up to equiv.

Ex: In our examples, we will have

| | \mathcal{W} | Cof | Fib | Final |
|---------------|--|--------------------|-----------------------------------|---------------|
| Top | \mathcal{W}_\sim -iso pair pairs | Some fibrations | Some cofibrations "all op." | all |
| sSet | \mathcal{W}_\sim -iso inj. inclusions | Kan fibrations | all | Kan cofibs |
| Ch^+ | \mathcal{W}_\sim -iso (proj=iso) | labeled maps | labeled maps | all |



2 Axioms for a model cat

Def: $f: X \rightarrow Y$ is a retract of $g: X' \rightarrow Y'$ if there is a diagram

$$\begin{array}{ccc} X & \xrightarrow{g} & X' \\ \downarrow f & \nearrow i & \downarrow g' \\ Y & \xrightarrow{h} & Y' \end{array}$$

s.t. $i \circ a = \text{id}_X$ and $b \circ id_Y = id_Y$

Def: A functional factorisation is a pair of functors

$d, p: \text{Mor}(\mathcal{C}) \rightarrow \text{Ho}(\mathcal{C})$

such that for $f: X \rightarrow Y$, we have

$$\begin{array}{ccc} X & \xrightarrow{d} & Y \\ \downarrow f & \nearrow p & \downarrow g \\ d(f) & \xrightarrow{p} & g \end{array}$$

Def: (Lifting property) Consider

$i: A \hookrightarrow B$

$p: X \rightarrow Y$

for each square, the following lift exists

$$\begin{array}{ccc} A & \xrightarrow{i} & B \\ \downarrow f & \nearrow p & \downarrow g \\ X & \xrightarrow{f} & Y \end{array}$$

then say that:

• i LLP wrt p

• p RLP wrt i

Ex: If \mathcal{C} has initial obj. 0 , then

B projective $\iff p: B \rightarrow 0$ has LLP wrt each epi $X \rightarrow 0$

Def: (Model categories) A model structure on \mathcal{C} is:

• Three full subcategories of $\text{Mor}(\mathcal{C})$

$\text{W}, \text{Cof}, \text{Fib}$

• Two functional factorisations (d, p) and (d', p') .

such that the following holds:

① Call further

$\mathcal{W} \cap \text{Cof}$ "weak cofibrations"

$\mathcal{W} \cap \text{Fib}$ "weak fibrations"

Then for each $f: X \rightarrow Y$, we have

$\alpha(f)$ cofibr. $\alpha'(f)$ fib. cofibr.

$\beta(f)$ fib. fib. $\beta'(f)$ fib. fib.

Def: (2-out-of-3) Consider

$X \xrightarrow{g} Z$

$Z \xrightarrow{h} Y$

Then it 2 out of $\{f, g, h\}$ are weak equivalences, so is the 3rd.

③ $\text{W}, \text{Cof}, \text{Fib}$ are closed under retracts.

④ Lifting: $i: A \hookrightarrow B$

$p: X \rightarrow Y$

for each square, the following lift exists

$$\begin{array}{ccc} A & \xrightarrow{i} & B \\ \downarrow f & \nearrow p & \downarrow g \\ X & \xrightarrow{f} & Y \end{array}$$

Thus, we get a cofibrant replacement

$Q: \mathcal{C} \rightarrow \mathcal{C}$

$R: \mathcal{C} \rightarrow \mathcal{C}$

together with a nat. eq.

$\beta: Q \circ i \rightarrow i \circ R$

$\alpha: Q \circ j \rightarrow j \circ R$

Def: Have restricted functors

$R: \mathcal{C}_\text{Cof} \rightarrow \mathcal{C}_\text{Cof}$

$Q: \mathcal{C}_\text{Fib} \rightarrow \mathcal{C}_\text{Fib}$

$\mathcal{W}: \mathcal{C} \rightarrow \mathcal{C}$

with $\mathcal{W} \cap \text{Cof} = \mathcal{C}_\text{Cof}$, $\mathcal{W} \cap \text{Fib} = \mathcal{C}_\text{Fib}$.

Ex: Let \mathcal{C} be a category. Then

\mathcal{C} cofib. $\iff \mathcal{C}$ LLP wrt. triv. fib.

Def: Only shows " \Leftarrow " of black version.

Can verify $f = p \circ i$ with

• i cofibr.

• p trivial fibration.

Then f LLP wrt. p

$\Rightarrow f$ retract of i

$\Rightarrow f$ cofibration. \square

Def: (Axiom 4) Let \mathcal{C} model cat. Then

\mathcal{C} cofib. $\iff \mathcal{C}$ LLP wrt. triv. fib.

Def: (Axiom 5) Let \mathcal{C} model cat. Then

\mathcal{C} LLP wrt. p \iff f retract of i

Def: (Axiom 6) Let \mathcal{C} model cat. Then

\mathcal{C} LLP wrt. i \iff i LLP wrt. p

Def: (Axiom 7) Let \mathcal{C} model cat. Then

\mathcal{C} LLP wrt. i \iff i LLP wrt. p

Def: (Axiom 8) Let \mathcal{C} model cat. Then

\mathcal{C} LLP wrt. p \iff p LLP wrt. i

Def: (Axiom 9) Let \mathcal{C} model cat. Then

\mathcal{C} LLP wrt. i \iff i LLP wrt. p

Def: (Axiom 10) Let \mathcal{C} model cat. Then

\mathcal{C} LLP wrt. p \iff p LLP wrt. i

Def: (Axiom 11) Let \mathcal{C} model cat. Then

\mathcal{C} LLP wrt. i \iff i LLP wrt. p

Def: (Axiom 12) Let \mathcal{C} model cat. Then

\mathcal{C} LLP wrt. p \iff p LLP wrt. i

Def: (Axiom 13) Let \mathcal{C} model cat. Then

\mathcal{C} LLP wrt. i \iff i LLP wrt. p

Def: (Axiom 14) Let \mathcal{C} model cat. Then

\mathcal{C} LLP wrt. p \iff p LLP wrt. i

Def: (Axiom 15) Let \mathcal{C} model cat. Then

\mathcal{C} LLP wrt. i \iff i LLP wrt. p

Def: (Axiom 16) Let \mathcal{C} model cat. Then

\mathcal{C} LLP wrt. p \iff p LLP wrt. i

Def: (Axiom 17) Let \mathcal{C} model cat. Then

\mathcal{C} LLP wrt. i \iff i LLP wrt. p

Def: (Axiom 18) Let \mathcal{C} model cat. Then

\mathcal{C} LLP wrt. p \iff p LLP wrt. i

Def: (Axiom 19) Let \mathcal{C} model cat. Then

\mathcal{C} LLP wrt. i \iff i LLP wrt. p

Def: (Axiom 20) Let \mathcal{C} model cat. Then

\mathcal{C} LLP wrt. p \iff p LLP wrt. i

Def: (Axiom 21) Let \mathcal{C} model cat. Then

\mathcal{C} LLP wrt. i \iff i LLP wrt. p

Def: (Axiom 22) Let \mathcal{C} model cat. Then

\mathcal{C} LLP wrt. p \iff p LLP wrt. i

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